

Exercise 1:  $\mathcal{L} : \mathbb{R}^2 \rightarrow \mathbb{R}$   
 $u(x_1, x_2) = \sqrt{x_1} + \sqrt{x_2}$ .

d. Show  $\mathcal{L}$  continuous, monotone, and convex

\*  $\mathcal{L}$  continuous iff it is preserved under taking limits

$$\left. \begin{array}{l} \forall \{x_n\}, \{y_n\} \quad x_n \succeq y_n \quad \forall n \\ x_n \rightarrow x \\ y_n \rightarrow y \end{array} \right\} \implies x \succeq y$$

equivalently iff  $\forall y, U(y)$  and  $L(y)$  are both closed.

$U(y) = \{x \mid u(x) \geq u(y)\}$  is closed because  $u$  is continuous.

$$x_n \succeq y_n \iff u(x_n) \geq u(y_n)$$

~~$u$  is continuous.~~  
 $x_n \rightarrow x$   
 $y_n \rightarrow y$

$x \succeq y$

$$\iff u(x) \geq u(y)$$

□

\*  $\succeq$  monotone.

iff.  $x \gg y \implies x \succ y$ .

ie.  $\forall i, x_i > y_i$

IMPACT  
IT IS EVEN  
STRONGLY  
MONOTONE.

Assume.  $x \gg y$ , ie  $\forall i, x_i > y_i$

$$\implies u(x) = \sqrt{x_1} + \sqrt{x_2} > \sqrt{y_1} + \sqrt{y_2} = u(y)$$

because  $x \mapsto \sqrt{x}$  is strictly increasing.

so  $u(x) > u(y)$ .

$\iff x \succ y \quad \square$

\*  $\succeq$  convex .  $\succeq$  convex . iff.

$x \succeq y$  and  $z \succeq y$ .

$\implies \forall \alpha \in [0, 1], \alpha x + (1-\alpha)z \succeq y$

~~it is equivalent to saying that  $\forall y, U(y)$  is a convex set.~~

which is equivalent to saying that  $u$  is quasi concave.

let  $x \succeq y$  and  $z \succeq y$ .

so  $u(x) \geq u(y)$  and  $u(z) \geq u(y)$

let  $\lambda x + (1-\lambda)z$

$$u(\lambda x + (1-\lambda)z) = \sqrt{\lambda x_2 + (1-\lambda)z_2} + \sqrt{\lambda x_2 + (1-\lambda)z_2}$$

$\geq$

$$\implies \begin{cases} z_1 \geq y_1 \\ z_2 \geq y_2 \end{cases}$$

$$\implies \begin{cases} x_1 \geq y_1 \\ \text{and } x_2 \geq y_2 \end{cases}$$

$$\text{so } \lambda x_2 + (1-\lambda)z_2 \geq \lambda y_2 + (1-\lambda)y_2 = y_2$$

same for  $z_1$

and  $\sqrt{\cdot}$  is increasing, so

$$\implies \sqrt{\lambda x_2 + (1-\lambda)z_2} + \sqrt{\lambda x_1 + (1-\lambda)z_1} \geq \sqrt{y_2} + \sqrt{y_1} = u(y)$$

Alternatively, we could have shown  $u$  is quasiconcave.

$$H_u(x) = \begin{pmatrix} \frac{\partial^2 u}{\partial x_1^2} & \frac{\partial^2 u}{\partial x_1 \partial x_2} \\ \frac{\partial^2 u}{\partial x_1 \partial x_2} & \frac{\partial^2 u}{\partial x_2^2} \end{pmatrix}$$

$$= \begin{pmatrix} \frac{\partial}{\partial x_2} \left( \frac{1}{2x_2} \right) & 0 \\ 0 & \frac{\partial}{\partial x_2} \left( \frac{1}{2} x_2^{-\frac{1}{2}} \right) \end{pmatrix} = \begin{pmatrix} -\frac{1}{4} x_2^{-3/2} & 0 \\ 0 & -\frac{1}{4} x_2^{-3/2} \end{pmatrix}$$

$H_u(x)$  has strictly negative eigenvalues  $\forall x \gg 0$ .

$\implies H_u(x)$  is negative definite.

$\implies u$  is strictly concave.

$\implies u$  is strictly quasiconcave.

$\iff \mathcal{L}$  is strictly concave.

$\implies \mathcal{L}$  is concave

$\square$

2) let  $p \gg 0$ ,  $w > 0$ .

$u$  is continuous.  $\implies$   $\dots$

$u$  is quasi-concave and monotone.

$\implies$  The Kuhn-Tucker conditions are sufficient.

$u$  is continuously twice-differentiable.

(not just necessary but...).

~~maximize~~  
~~minimize~~  
 UMP  $\left\{ \begin{array}{l} \max u(x) \\ x \in X \\ p \cdot x \leq w. \end{array} \right.$

$\bar{x} \in x(p, w) \equiv$  Walrasian demand corresp.

of.  $\nabla u(\bar{x}) = \lambda p$   
 and  $\lambda(p \cdot \bar{x} - w) = 0$ .

$$\begin{pmatrix} \frac{\partial u}{\partial x_1}(\bar{x}) \\ \frac{\partial u}{\partial x_2}(\bar{x}) \end{pmatrix} = \begin{pmatrix} \frac{1}{2\sqrt{x_1}} \\ \frac{1}{2\sqrt{x_2}} \end{pmatrix} = \lambda \begin{pmatrix} p_1 \\ p_2 \end{pmatrix}$$

we could also argue  $u$  strictly concave  $\implies \lambda > 0$

$u$  concave  $\implies u$  locally max satisfied.

$\implies$  Kuhn-Tucker. i.e.,

$$\underline{\underline{p \cdot \bar{x} = w}}$$

So ...  $\lambda(p \cdot \bar{x}) = \lambda w$ .

$$= \frac{1}{2} (\sqrt{x_1} + \sqrt{x_2}) = \frac{1}{2} u(\bar{x}).$$

$$\implies \lambda = \frac{u(\bar{x})}{2w}$$

$$\implies \frac{1}{2\sqrt{x_i}} = \frac{u(\bar{x})}{2w} p_i$$

$$\frac{1}{\sum x_i} = \frac{u(x)}{\sum w p_i}$$

$$\Rightarrow x_i = \frac{w}{p_i \cdot u(x)}$$

$$u(x) = \frac{w}{u(x)} \left( \frac{1}{p_1} + \frac{1}{p_2} \right)$$

$$u(x)^2 = w \frac{p_1 + p_2}{p_1 p_2}$$

$$\Rightarrow \cancel{x_i = \frac{w}{p_i} \frac{p_1 + p_2}{p_1 p_2}}$$

OR!

$$x_i = \frac{w}{p_i} \frac{1}{w} \frac{p_1 p_2}{(p_1 + p_2)}$$

□

$$x_1 = \frac{w}{p_1 + p_2} \cdot \frac{p_2}{p_1} \quad \text{and} \quad x_2 = \frac{w}{p_1 + p_2} \cdot \frac{p_1}{p_2}$$

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## Exercise 2:

$$x_1 = \frac{w}{p_2}$$

$$x_2 = \frac{w(p_2 - p_1)}{p_2^2}$$

1) Hom. of deg 0. H.

$$x(tp, tw) = \begin{pmatrix} \frac{tw}{tp_2} \\ \frac{t^2 w (p_2 - p_1)}{t^2 p_2^2} \end{pmatrix} = \begin{pmatrix} \frac{w}{p_2} \\ \frac{w(p_2 - p_1)}{p_2^2} \end{pmatrix} = x(p, w)$$

So YES -

□

2) Walras' law.

$$p \cdot x = p_1 \frac{w}{p_2} + p_2 \frac{w(p_2 - p_1)}{p_2^2}$$

$$= w$$

So YES holds □



3) WARP. in the framework of the demand.

~~that~~ satisfies WARP iff.  
 $x(p, w)$ .

$$p' \cdot x(p, w) - w' \leq 0.$$

and.  $x(p', w') \neq x(p, w)$

$$\implies p \cdot x(p', w') - w > 0.$$

4) Set  $p_2 = 1$ .

$$x_2 = w.$$

$$x_2 = w(1 - p_1).$$

Suppose  $p' \cdot x(p, w) \leq w'$ .

i.e.  $p_1' w + 1 \cdot w \cdot (1 - p_1) \leq w'$ .

~~$$w(1 + p_1' - p_1) \leq w'$$~~

that it cannot have.

~~$$p \cdot x(p', w') = (p_1 - p_1' + 1) w'$$~~

The question is, can I find  
 $x \neq x'$  with

$p \cdot x' \leq w$  and  $p' \cdot x \leq w'$  ?

assume  $\begin{cases} p \cdot x' \leq w \\ p' \cdot x \leq w' \end{cases} \Leftrightarrow \begin{cases} w'(1 + p_2' - p_2) \leq w \\ w(1 + p_2 - p_2') \leq w' \end{cases}$



note:  $p_2 < 1$   $\Leftrightarrow$   $\begin{cases} w' - w \leq (p_2' - p_2)w' \\ w' - w \geq (p_2' - p_2)w \end{cases}$

let  $w = w' \Rightarrow p_2' = p_2$   
 does not work.

suppose  $w' = w + \Delta w$   
 $\Delta w > 0$

$0 \leq (p_2' - p_2) \Delta w$   
 $\Delta w > 0$   
 $\Rightarrow$

$\Delta w \leq (p_2' - p_2)(w + \Delta w)$  (1)  
 $\Delta w \geq (p_2' - p_2)w$  (2)

~~$\Delta w \leq (p_2' - p_2)w + \Delta w(p_2' - p_2)$~~

So take for instance.

$$p_2' = p_2 + \underbrace{\Delta p_2}_{\geq 0}$$

$$\Delta w \leq \Delta p_2 (w + \Delta w)$$

$$\Delta w \geq \Delta p_2 w$$



So long as.

$$\Delta p_2 \leq \frac{\Delta w}{w}$$

then we can find  
counter example to

WARP!

□

### Exercise 3:

1)  $\mathcal{L}$  satisfies the independence axiom.

$$\text{iff } \forall L, L', L'' \in \mathcal{L},$$

$$\forall d \in (0, 1). \quad \text{caused.}$$

$$L \sim L' \implies \underbrace{dL + (1-d)L''}_{\text{this is a compound lottery}} \sim dL' + (1-d)L''$$

"this is a compound lottery".

2)  $U: \mathcal{L} \rightarrow \mathbb{R}$  has the expected utility form.

iff  $\exists (u_1, \dots, u_n)$  assignment to each outcomes  $(c_1, \dots, c_n)$ .

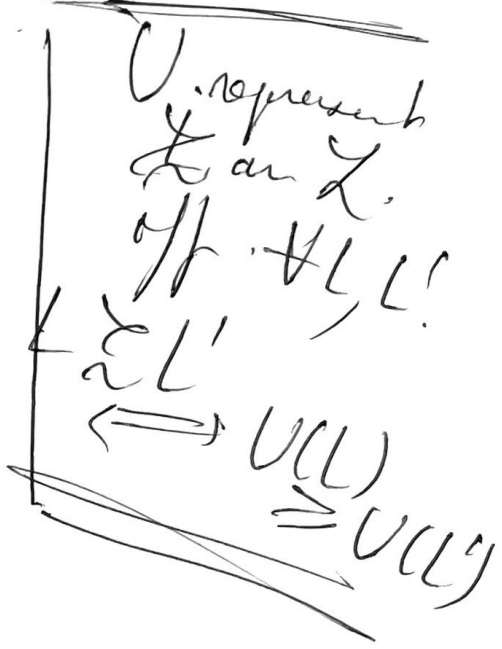
$$\text{with } \forall L = (p_1, \dots, p_n) \in \mathcal{L}$$

$$U(L) = \sum_{i=1}^n u_i p_i$$

2) h) Show that  $\Sigma$  satisfies the independence  
 axiom

a) first,

let  $L, L', L''$   
 $\overset{u}{(p_i)}$   $\overset{u}{(p_i')}$   $\overset{u}{(p_i'')}$



$\implies$  let  $d \in [0, 1]$ .



let  $L \sim L'$



$$U(L) = U(L')$$

$$U(\alpha L + (1-\alpha)L'')$$

$= \sum_i u_i(\alpha p_i + (1-\alpha)p_i'')$   
 the compound lottery is also representable.

$$= \alpha \sum_i u_i p_i + (1-\alpha) \sum_i p_i'' u_i$$

$$= \alpha U(L) + (1-\alpha) U(L'')$$

$$= \sum_i u_i (\alpha p_i + (1-\alpha)p_i'')$$

$$= \sum_i u_i (\alpha p_i + (1-\alpha)p_i'') = U(\alpha L + (1-\alpha)L'')$$

INDEPENDENCE AXIOM

