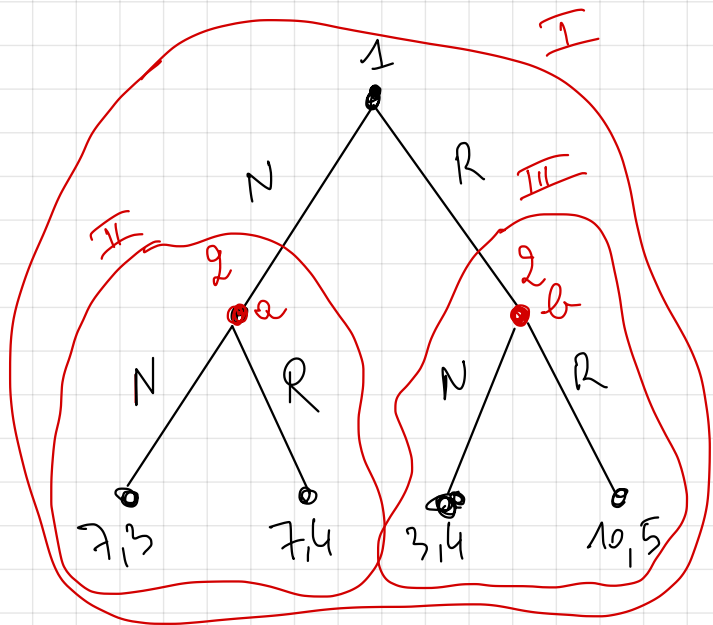


course #9  
March 26th.  
2026.





two players: 1 and 2

Set of strategies  $S_1 = \{N, R\}$

$S_2 = \{NN, RR, NR, RN\}$

player 2

	NN	RR	NR	RN
N	7,3	7,4*	7,3	7,4*
R	3,4	10,5*	10,5*	3,4

player 1

A strategy for player 2 consists in choosing an action for each node where she has to play.

a	b
N	R

Subgame perfect Nash equilibrium  $\rightarrow$  it has to be a N.E

in each subgame, i.e. including the whole game. (I)

SPNE: (R, RR).  $\rightarrow$  payoffs (10, 5).

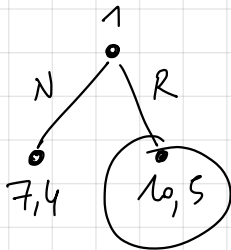
Backward induction.

NE: (R, NR) SP? No. because N is not optimal at node 2 (subgame II) for player 2.

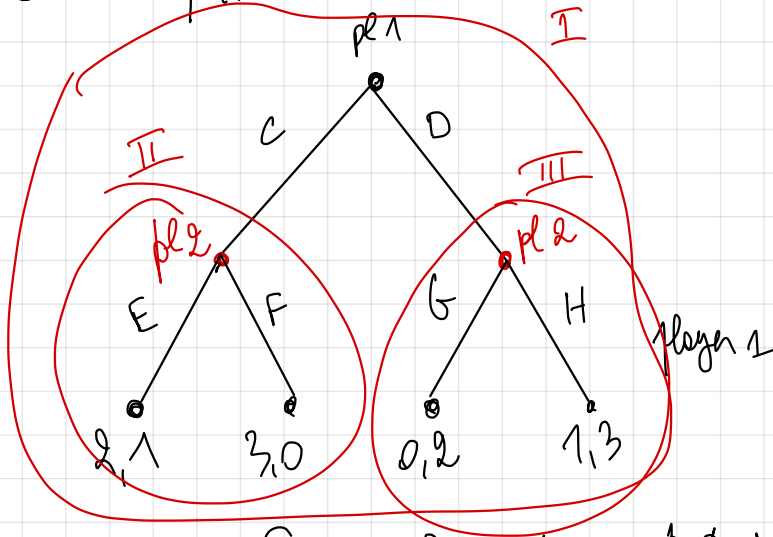
Subgame II:  $P_2 \rightarrow R \rightarrow 7, 4$

III:  $P_2 \rightarrow R \rightarrow 10, 5$

Now, work with the following reduced game:



osborne p 170.



$$S_1 = \{C, D\}$$

$$S_2 = \{EG, EH, FG, FH\}$$

player 2

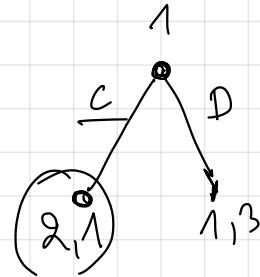
	EG	EH	FG	FH
C	2, 1*	2, 1*	3, 0	3, 0
D	0, 2	1, 3*	0, 2	1, 3*

Two N.E in pure strategies: (C, EG) and (C, EH).

SPNE: backward induction

II: pl 2  $\rightarrow$  E  $\rightarrow$  2, 1

III: pl 2  $\rightarrow$  H  $\rightarrow$  1, 3



# CHAPTER 6: Extensive games with perfect information

## Illustrations

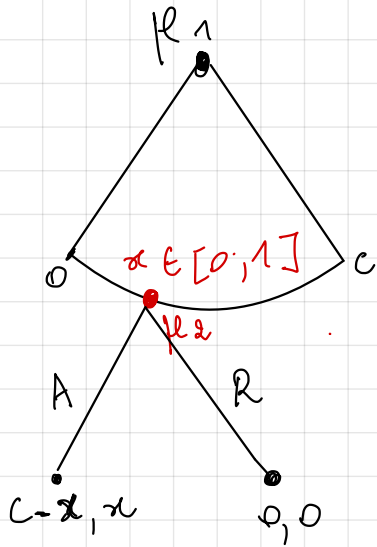
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### The ultimatum game

- 2 players: 1 and 2
- bargaining over the division of a "pie" let's say  $c \in (100 \text{ €})$ .
  - 1) player 1 proposes a split  $x$  for player 2,  $c-x$  for player 1.
  - 2) player 2 accepts or rejects.
- final payoffs:  $(c-x, x)$  or  $(0, 0)$ .
- terminal histories  $(x, z)$  where

$x \in [0; c]$        $c \in \mathbb{R}^+$ .

$Z \in \{ \text{Accept}, \text{reject} \}$



• Player function  $P(\phi) = P_1$   
 $P(x) = P_2$

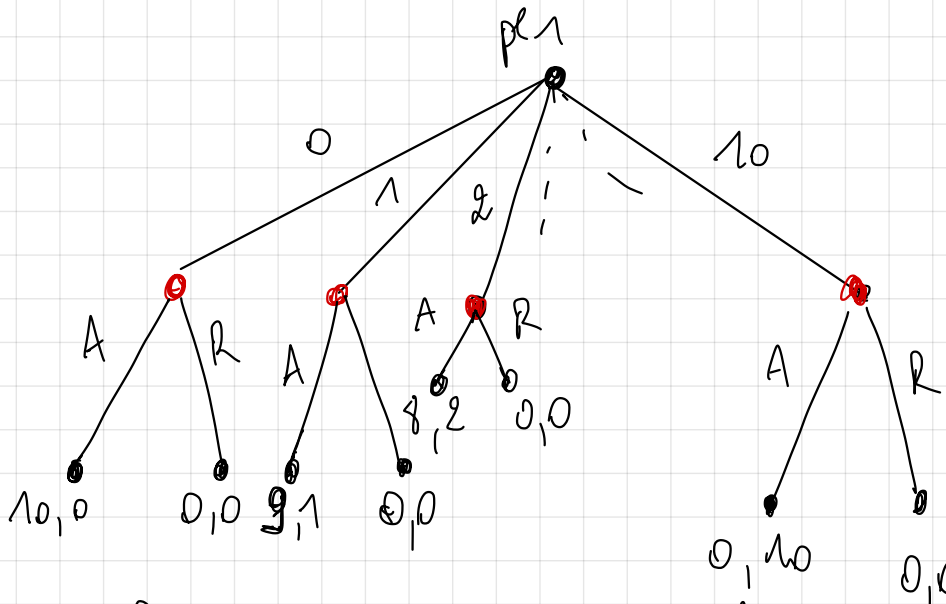
•  $u_1(x, \text{Accept}) = c - x$

•  $u_2(x, \text{Accept}) = x$

•  $u_i(x, \text{Reject}) = 0$   
 $i = 1, 2$

Discrete version :

$C = 10$



What are the subgame perfect NE?

Backward induction, look at subgames of length 1:

at each such subgame, Player 2 faces  $x \in [0; 10]$ .

- If  $x > 0$  : she accepts.
- If  $x = 0$  : she is indifferent between R or A.

In a SPNE, player 2 accepts everything (including 0)  
or accepts  $x > 0$  and rejects  $x = 0$ .

whole game, optimal strategy for player 1.

• If player 2 accepts everything, then it's optimal for pl. 1 to propose  $x = 0$ .  $\rightarrow$  final payoff 1, 0.

• If player 2 accepts  $x > 0$  and rejects  $x = 0$ , then pl. 1 proposes  $x = 1$ .  $\rightarrow$  final payoff 0, 1.

Back to the continuous game  $C \in \mathbb{R}^+$ , and  $x \in [0; C]$ .

For player 2, it does not change anything:

Player 1: • if player 2 accepts everything, player 1 proposes  $x = 0$   
• if player 2 accepts  $x > 0$  and rejects  $x = 0$ , then does

not exist an optimal strategy for player 1.

The only <sup>weak</sup> SPNE :  $\alpha = 0$ , player 2 accepts everything.

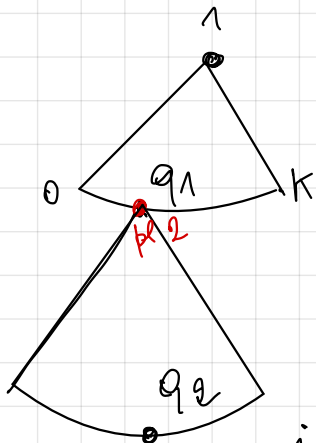
Experiments: 70's.

- $\alpha \approx 65$  or  $70\%$  of the amount.
  - about 20% of the offers were rejected.
- Equity. 50/50.
- risk aversion / uncertainty.

Nilton Friedman : good model , hypothesis may be unrealistic as long as the prediction is good.

Stackelberg Model : sequential Cournot.

2 firms , 1 is the leader , and 2 is the follower.  
firm 1 set its quantity first :  $q_1$ . Then firm 2 sets  $q_2$ .



• Firm 1 moves first ( $q_1, q_2$ ).

• payoffs :  $Q = q_1 + q_2$

$P(Q)$  = inverse demand

$C_i(q_i)$  : individual total cost

$$i = 1, 2 : \pi_i(q_i, q_{-i}) = q_i P(Q) - C_i(q_i).$$

Simplify calculations:  $P(Q) = 2 - Q$   $C_i(q_i) = q_i$   
 $= 2 - q_1 - q_2$  (constant marginal cost = 1).

SPNE, backward induction quantity  $\times$  price - total cost.

$$\begin{aligned}\text{Firm 2: } \max \pi_2 &= q_2 (2 - q_1 - q_2) - q_2 \\ &= 2q_2 - q_1 q_2 - q_2^2 - q_2 \\ &= q_2 - q_1 q_2 - q_2^2\end{aligned}$$

$$\frac{\partial \pi_2}{\partial q_2} = 1 - q_1 - 2q_2 = 0 \Leftrightarrow q_2^* = \frac{1 - q_1}{2}$$

Now, we look at the first stage, player 1 has to choose  $q_1$  knowing that  $q_2 = \frac{1 - q_1}{2}$ .

Firm 1's problem:

$$\max \pi_1 = q_1 (2 - q_1 - q_2) - q_1$$

$$\begin{aligned}
 &= q_1 \left( 2 - q_1 - \left( \frac{1 - q_1}{2} \right) \right) - q_1 \\
 &= 2q_1 - q_1^2 - q_1/2 + q_1^2/2 - q_1 \\
 &= \frac{q_1}{2} - \frac{q_1^2}{2}
 \end{aligned}$$

$$\frac{\partial \pi_1}{\partial q_1} = \frac{1}{2} - \frac{2q_1}{2} = 0 \Leftrightarrow q_1^* = \frac{1}{2}$$

$$\begin{aligned}
 \text{So, } q_2^* &= \frac{1 - q_1}{2} \\
 &= \frac{1 - 1/2}{2} = \frac{1}{4}
 \end{aligned}$$

SPNE  $\left( \frac{1}{2}; \frac{1}{4} \right)$ .  
 $\rightarrow \pi_1 > \pi_2$ .

Next week: ch 9 Bayesian games -

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